UNIVERSITÄT DUISBURG ESSEN

Internet-Technologie & Web Engineering
Introduction

Dr.-Ing. Matthäus Wander

Universität Duisburg-Essen

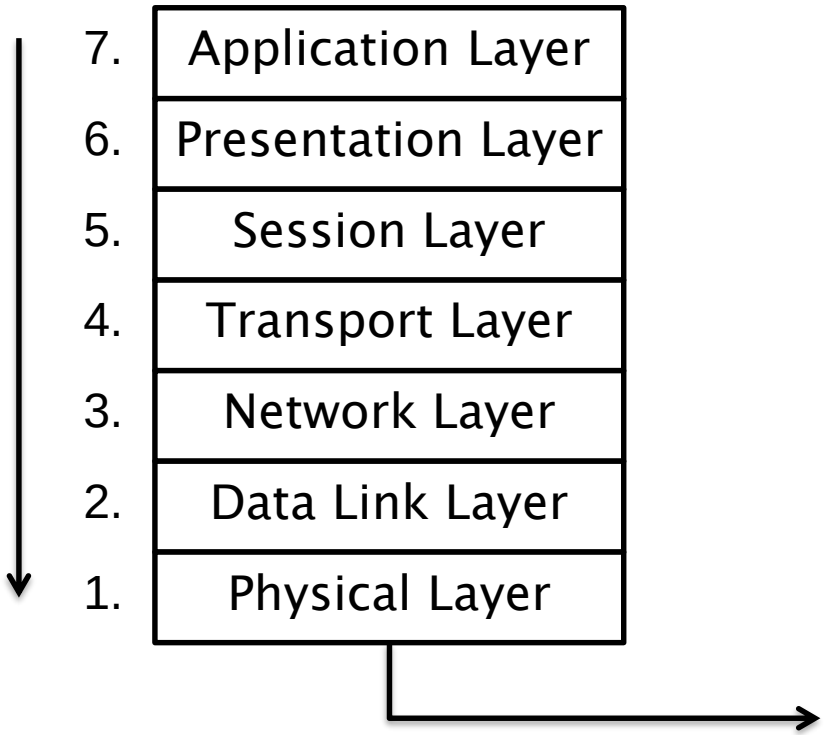
Goal

- Applications run on more than one computer
- Why?
  - Communicate with people (Facebook, Skype)
  - Share files (Dropbox, BitTorrent)
  - Access remote data/functionality (LSF, online shop)

⇨ **Distributed Application** or **Distributed System**

- What do we need?
  - Computer network
  - Rules for network communication

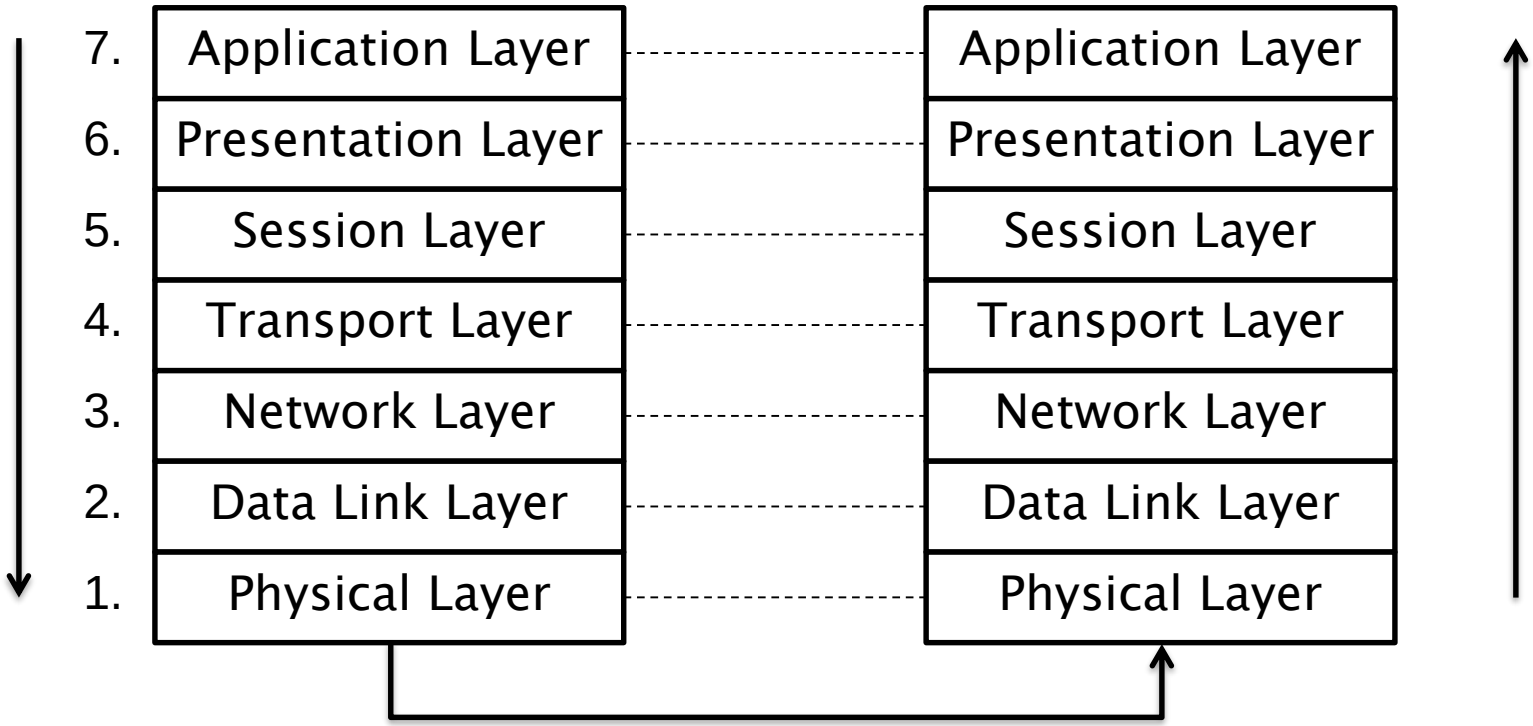
ISO/OSI Network Model



Sender:

- Protocol Data Unit (PDU) handed to bottom layer
- Each layer prepends PDU from upper layer with a header (or footer)

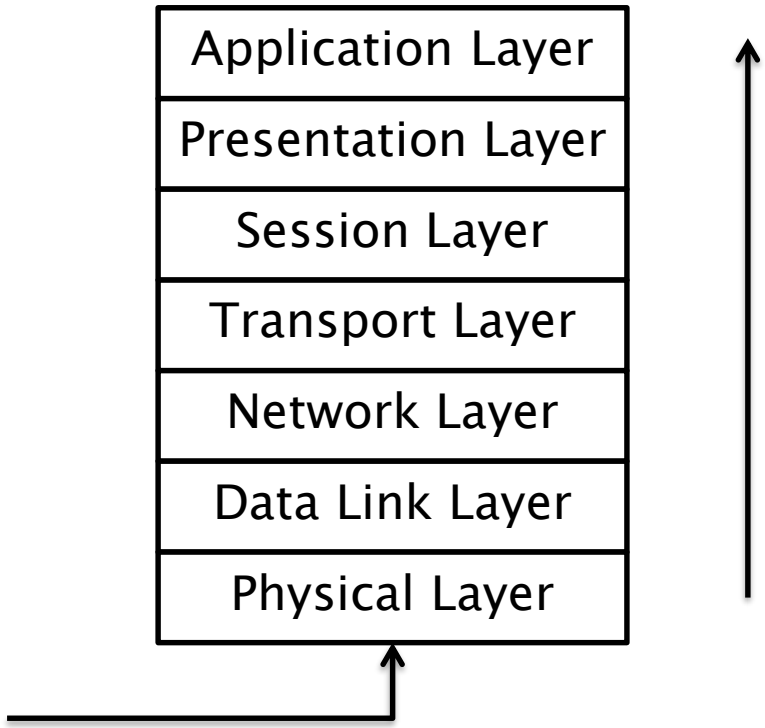
ISO/OSI Network Model



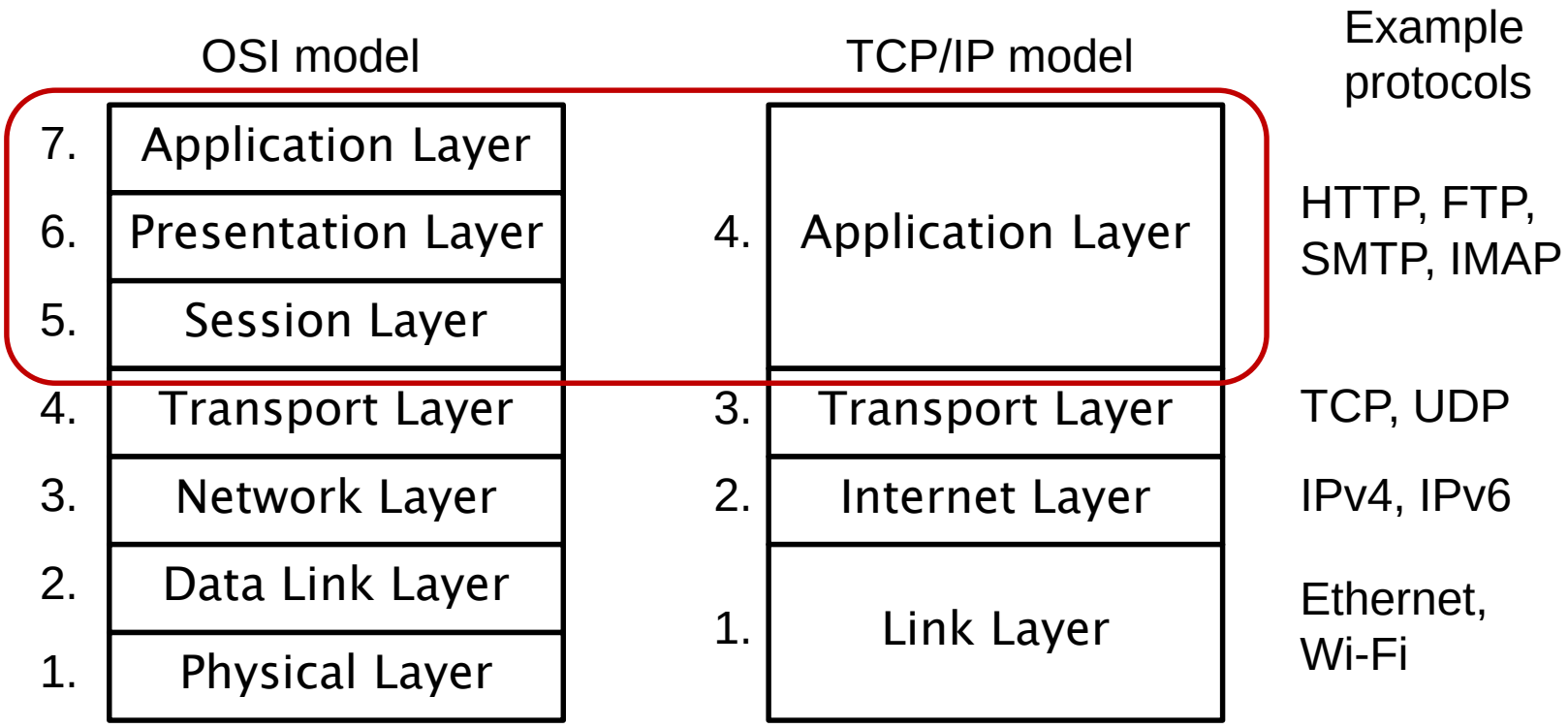
ISO/OSI Network Model

Receiver:

- Each layer strips header (or footer) from PDU, performs its function
- Gives PDU to upper layer



Comparison: OSI Model vs. TCP/IP Model



Link Layer – Physical and Data Link Layer

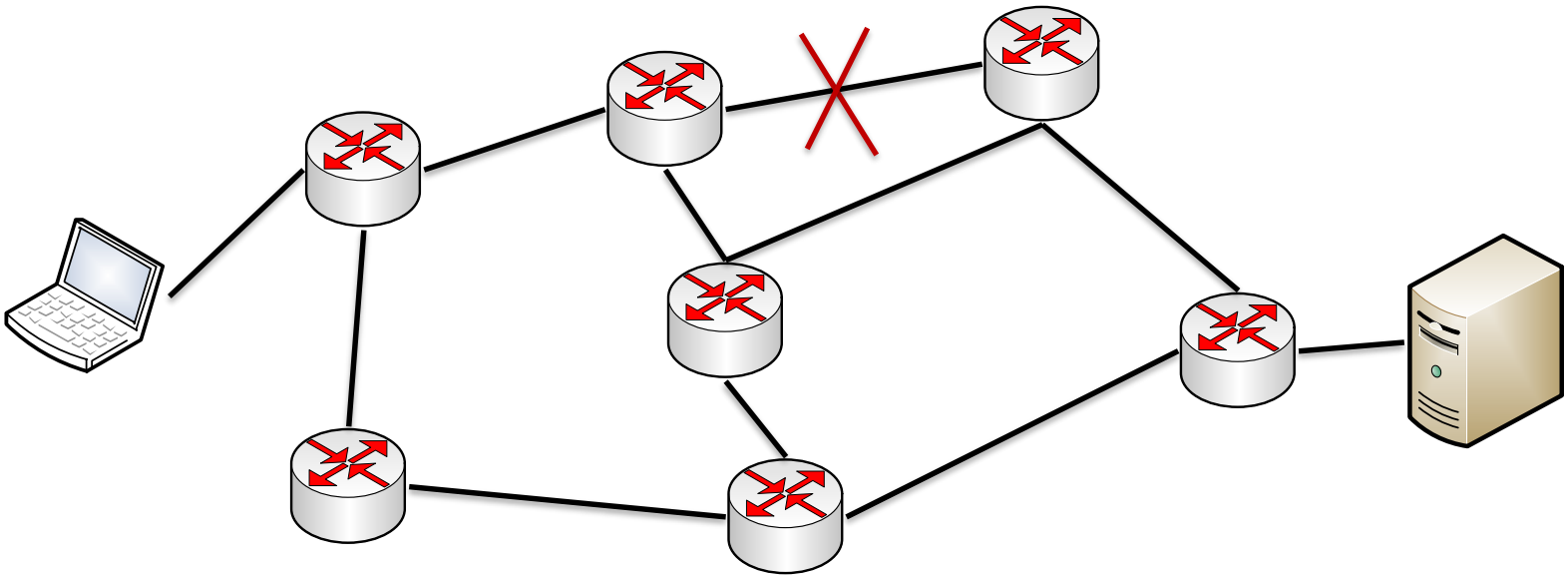
- Communication interface to the local network
  - Data transmission of directly connected computers
- Protocols: Ethernet, Wi-Fi, Bluetooth, ZigBee, …
- Data units: **frames**
- Adress type: MAC address (48/64 bit)
  - Example: 00:80:41:ae:fd:7e
- Tasks:
  - Cooperative access to (shared) network media
  - Bit encoding

Internet Layer – Network Layer

- Data transmission between computers (**hosts**)
- Protocols: IPv4, IPv6, …
- Data unit: **packets**
- Address type: IP address (32/128 bit)
  - Examples: 134.91.78.133 2001:638:501:8efc::133
- Tasks:
  - Routing: find best way for packet through network
  - Fragmentation: split large packet into smaller packets
  - Network congestion: handle overloaded network links

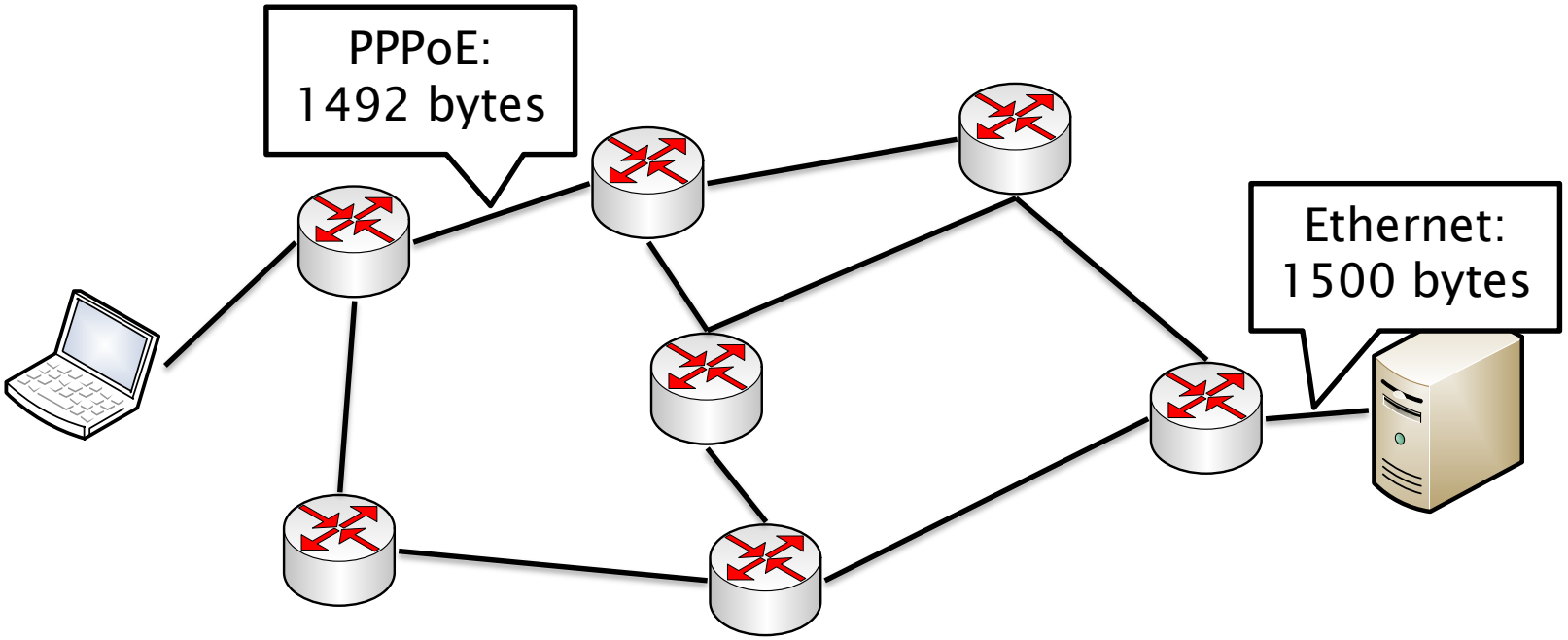
Routing

- Find best way for packet through network
  - Within a network (e.g. campus) & between networks
- Keep track of network topology when network links fail



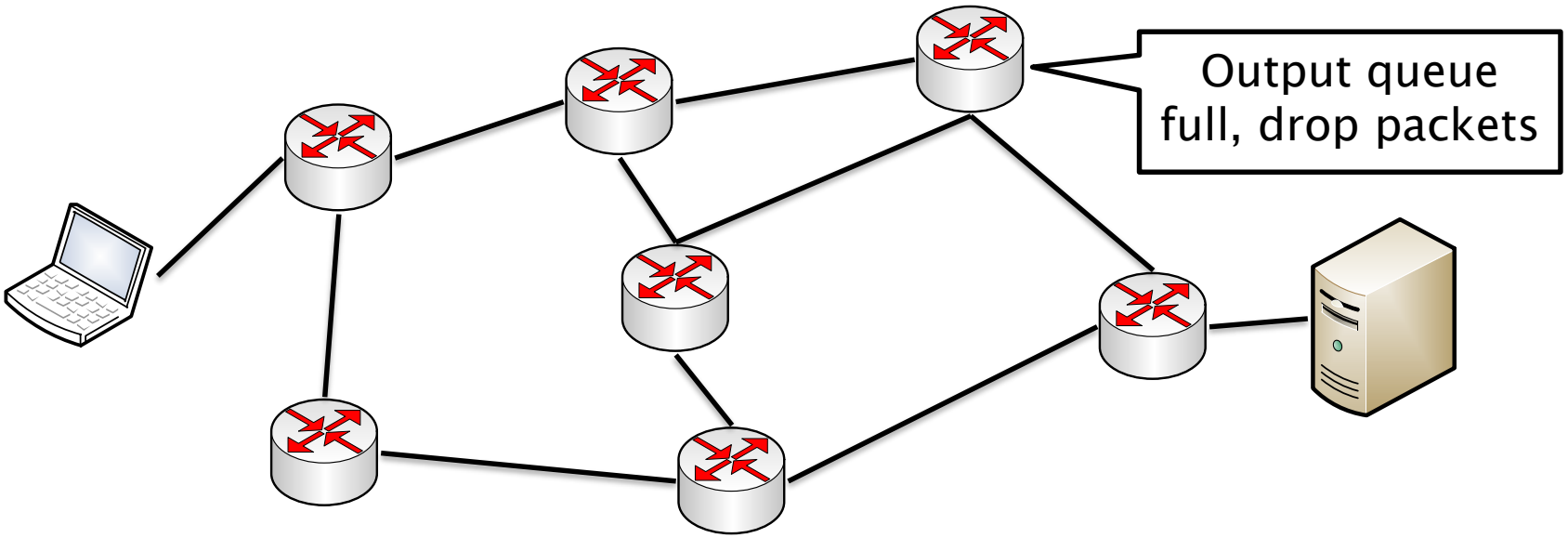
Fragmentation

- Split large packet into smaller packets
- Network links have packet size limitations
  - **Maximum Transmission Unit** (MTU)



Network Congestion

- Handle overloaded network links
- Drop packet silently (other layer must detect packet loss and retransmit)
  - Or send *Explicit Congestion Notification* (ECN)



Transport Layer

- Data transmission between processes
- Protocols: TCP, UDP, …
- Data unit: **segment** (TCP), **datagram** (UDP)
- Address type: port number (16 bit)
  - Examples: 80, 443, 51539
- Tasks:
  - Deliver reliable byte stream between processes (TCP)
  - Deliver individual messages with low latency (UDP)

Transmission Control Protocol (TCP)

- Connection-oriented
  - Connection establishment and termination
  - Delivers a continuous byte stream on top of packet-switched network
- Congestion control: determine transmission rate
- Detects and handles network errors
  - Lost, duplicate or out-of-order packets
  - Acknowledgement & retransmission, re-ordering
- Suitable for reliable transmissions and large data amounts

User Datagram Protocol (UDP)

- Connectionless
  - Best-effort attempt to deliver a datagram in one packet
  - Efficient due to low overhead/functionality
- Unreliable
  - Detects truncated/altered datagrams with checksum
  - No other error handling, no retransmission
- Suitable for low latency applications that handle errors themselves (e.g. VoIP, some online games)

Application Layer

- Application-specific data transmission and processing
  - Web: obtain document from web server via HTTP
  - Email: transfer email to another mailbox via SMTP
- Protocol, data unit, addressing and tasks different for each and every application
  - Web: client/server system with requests and responses
  - BitTorrent: peer-to-peer system with messages
  - VoIP: audio stream on top of messages

Network Protocol

- Challenge: **interoperability**
  - Different computer hardware and operating systems
  - Different software implementations and vendors
  - Different feature sets and extensions
- **Network protocol**: rules for interaction
  - Defines what to send when, and what it means
  - Syntax: message format (which bytes to send)
  - Semantics: meaning of messages and bytes
  - Clarifies the communication, but not application design choices (e.g. user interface, local file storage)

Internet Standardization

- Who publishes Internet standards?
- IEEE: communication technology
  - Ethernet, Wi-Fi, Bluetooth, …
- IETF: Internet protocols
  - IP, TCP, HTTP, DNS, FTP, SMTP, …
- W3C: World Wide Web standards
  - HTML, CSS, XML, SVG, …
- And various others, e.g. ECMA (JavaScript, JSON), ISO/IEC (JPEG, MP3), ITU (H.323)

Internet Engineering Task Force (IETF)

"We believe in: rough consensus and running code."
– David D. Clark, 1992

- IETF is an open standardization organization
  - Run by volunteers (usually with jobs in industry)
  - Anyone can participate (no membership required)
- Coordination of protocol engineering
  - Working groups, mailing lists, international meetings
  - Public resources: https://www.ietf.org

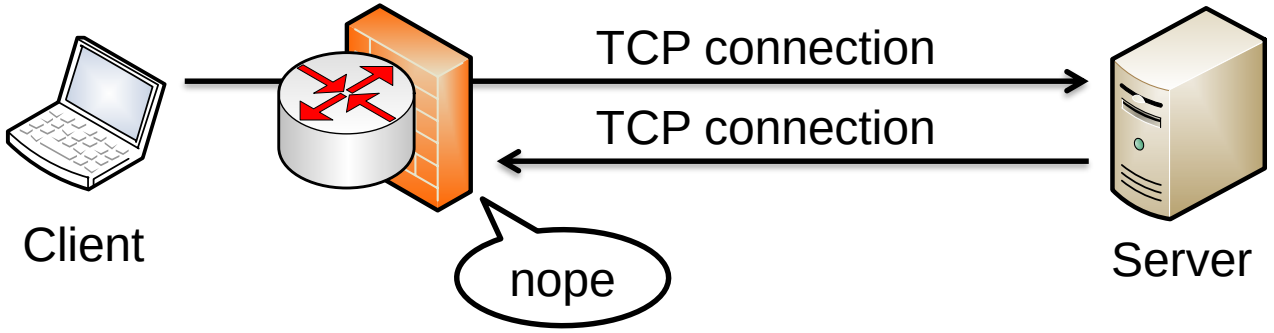
IETF Publications

- Request for Comments (RFC)
  - e.g. RFC 768: *User Datagram Protocol*
  - Various types of RFCs: (Proposed) Standard, Informational, Experimental, Obsolete
- Internet Standard (STD)
  - Well-known, mature and stable specification
  - e.g. STD 6: *User Datagram Protocol* (same as RFC 768)
  - Process: Internet Draft (I-D) ⇨ Proposed Standard (RFC) ⇨ Internet Standard (STD)
  - Very few protocols become STD (e.g. HTTP is not)

IETF Publications

- Best Current Practice (BCP)
  - Usually operational advice (how to run networks)
- IETF publications use specific terminology
  - e.g. MUST, MUST NOT, SHOULD, SHOULD NOT, MAY
- RFCs never change
  - Corrections published separately as Errata
  - New RFCs may update or obsolete old RFCs
- RFCs do not always reflect state of the art
  - New publications take time and effort

Model of Internet Communication



- Client host connected to the Internet
  - Can create outgoing connections
  - Local network does not allow incoming connections
  - Port forwarding necessary (manually, or via UPnP)
- Server host allows connections in and out

Further Challenges of Internet Applications

- Parallel activities
  - Autonomous components executing concurrent tasks
  - One server deals with >1 clients (do not block!)
- Communication via message passing
  - No shared memory, but network delays
- No knowledge of global state or global clock
  - Each client/server has their limited point of view
- No absolute trust in the other side
  - Assume buggy software or a malicious attacker