

# GPU-based NSEC3 Hash Breaking

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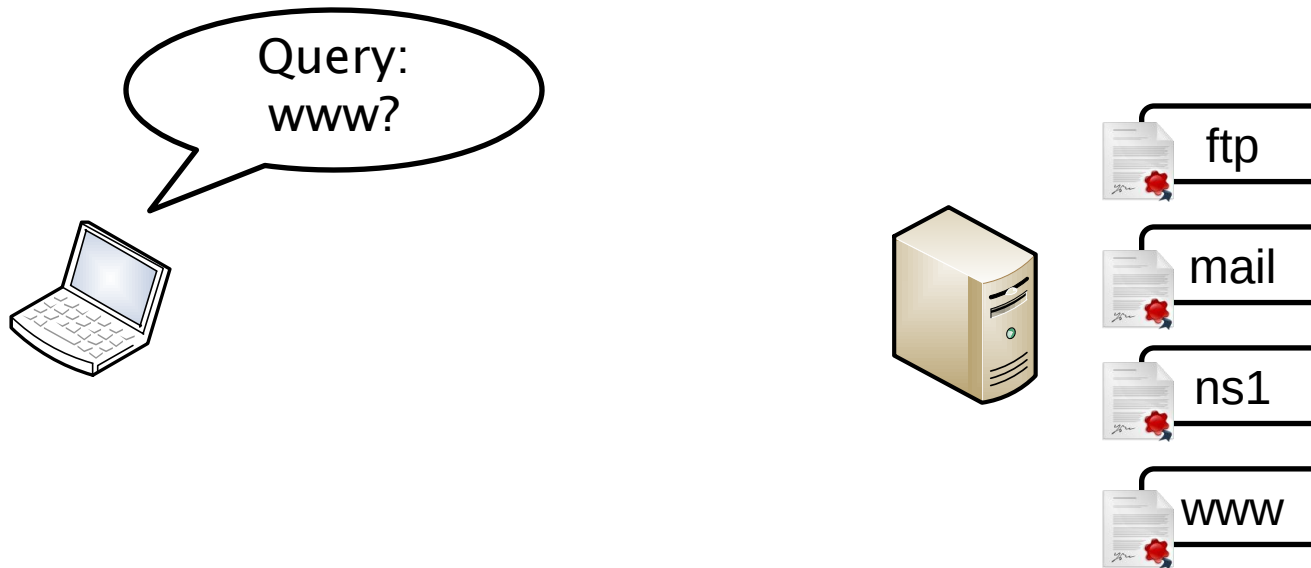
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IEEE NCA 2014

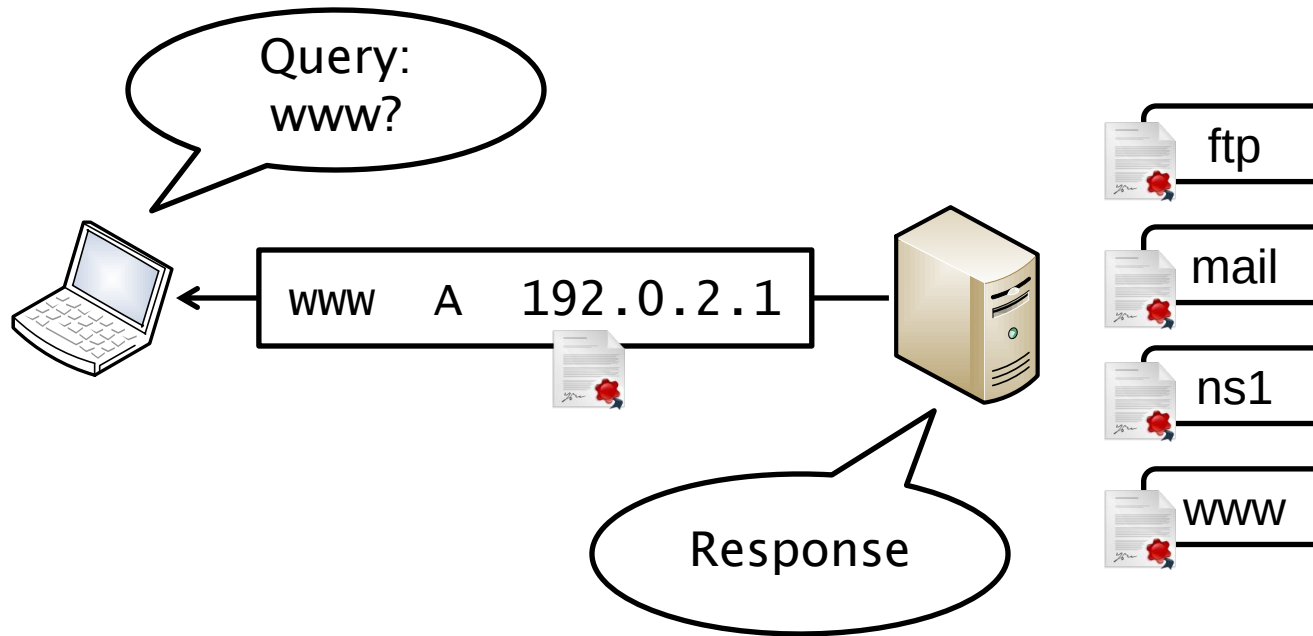
Cambridge, August 22, 2014

# NSEC3 FUNDAMENTALS

# Secure Denial of Existence with DNSSEC

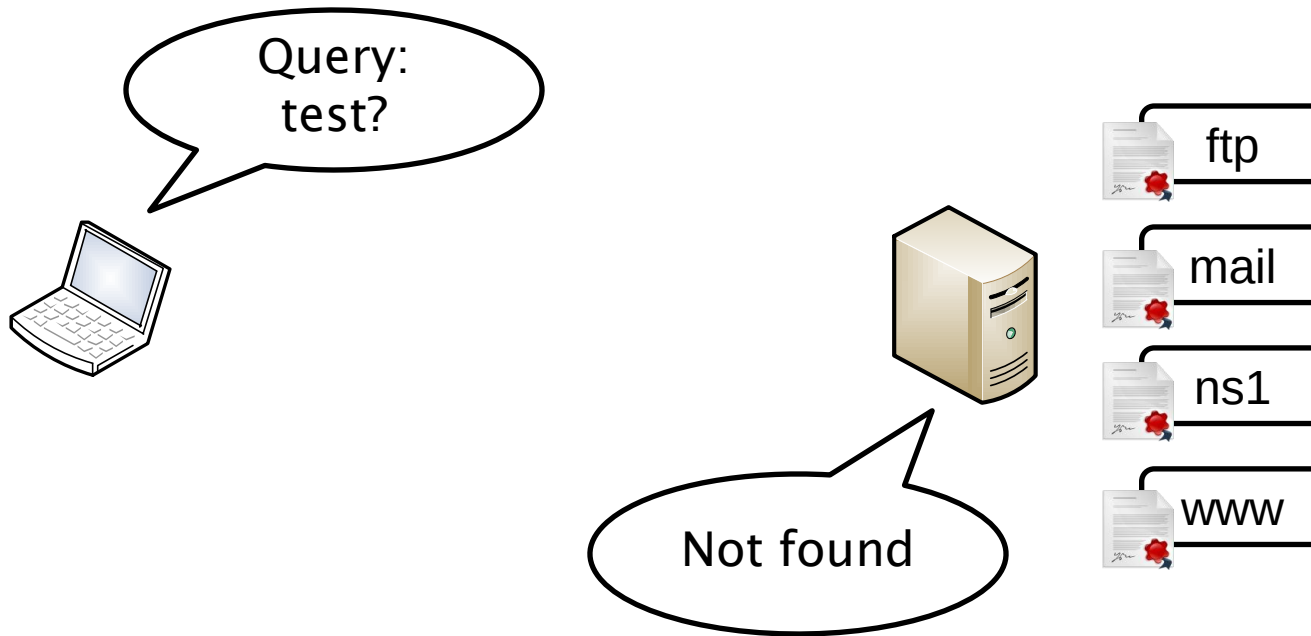


# Secure Denial of Existence with DNSSEC



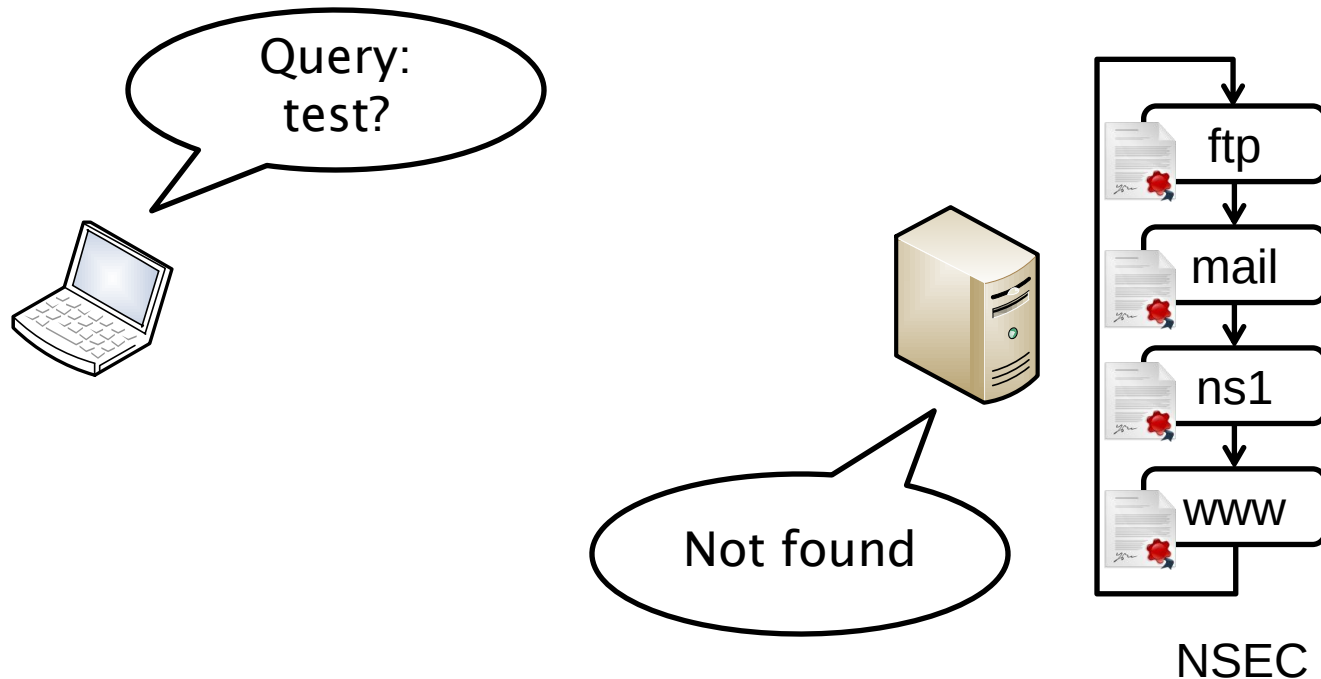
- Signed resource record
- Authenticity verified with PKI built into DNSSEC

# Secure Denial of Existence with DNSSEC



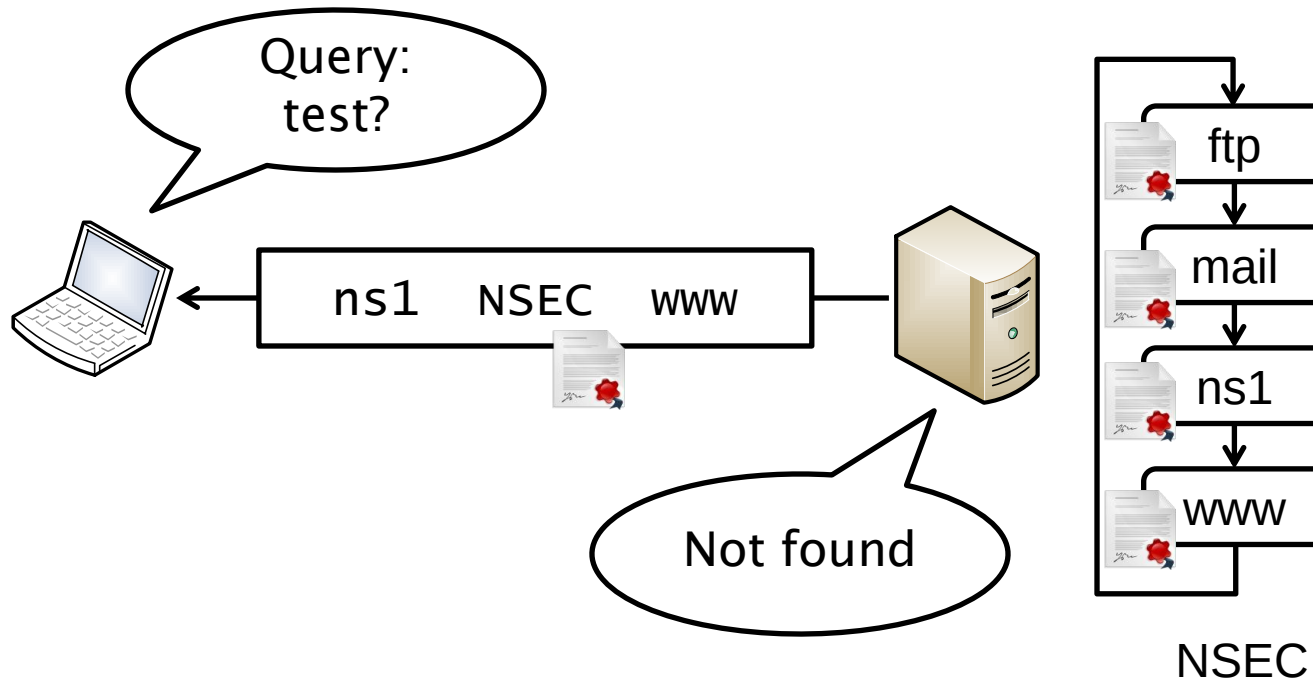
- Query name:  $n$

# Secure Denial of Existence with DNSSEC



- Query name:  $n$

# Secure Denial of Existence with DNSSEC



- Query name:  $n$
- Response:  $(r_1, r_2)$  with  $r_1 < n < r_2$

# Zone Enumeration

- Copy DNSSEC zone by crawling NSEC records
- Nominet (.uk) position:

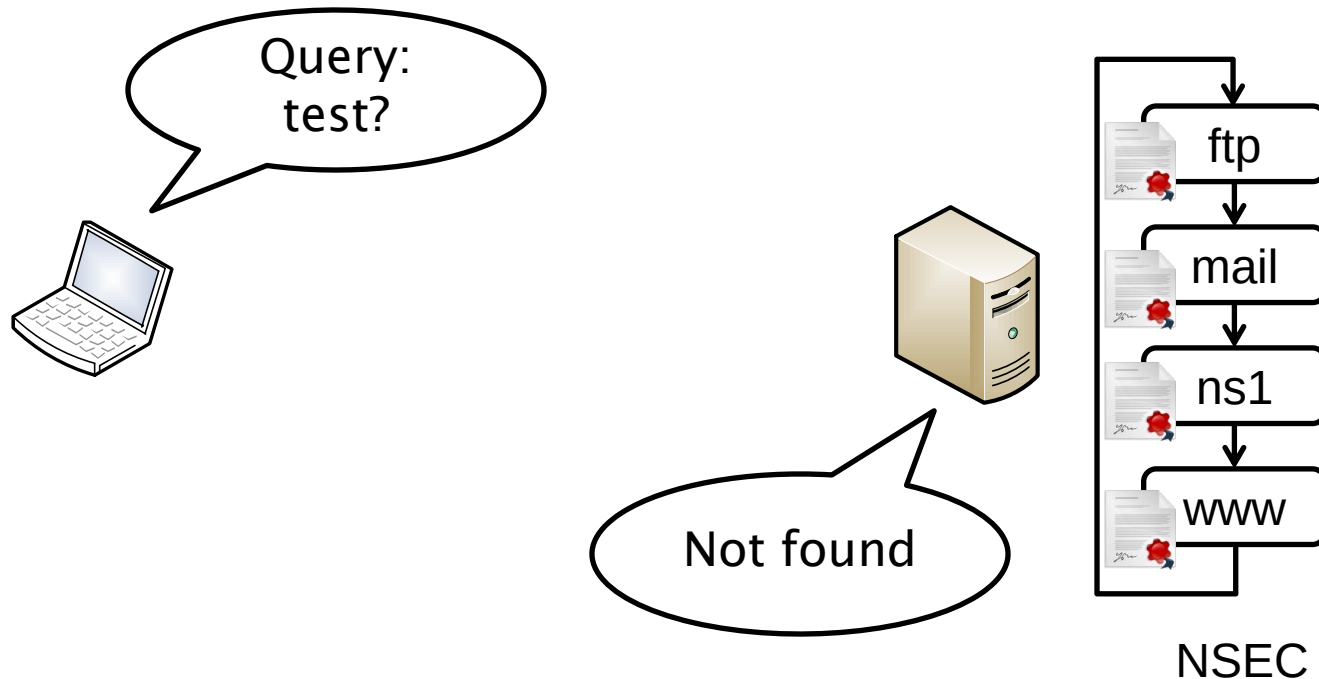
„Zone file enumeration is a show-stopper for us that will prevent us from fully implementing DNSSEC.“

- DENIC (.de) position:

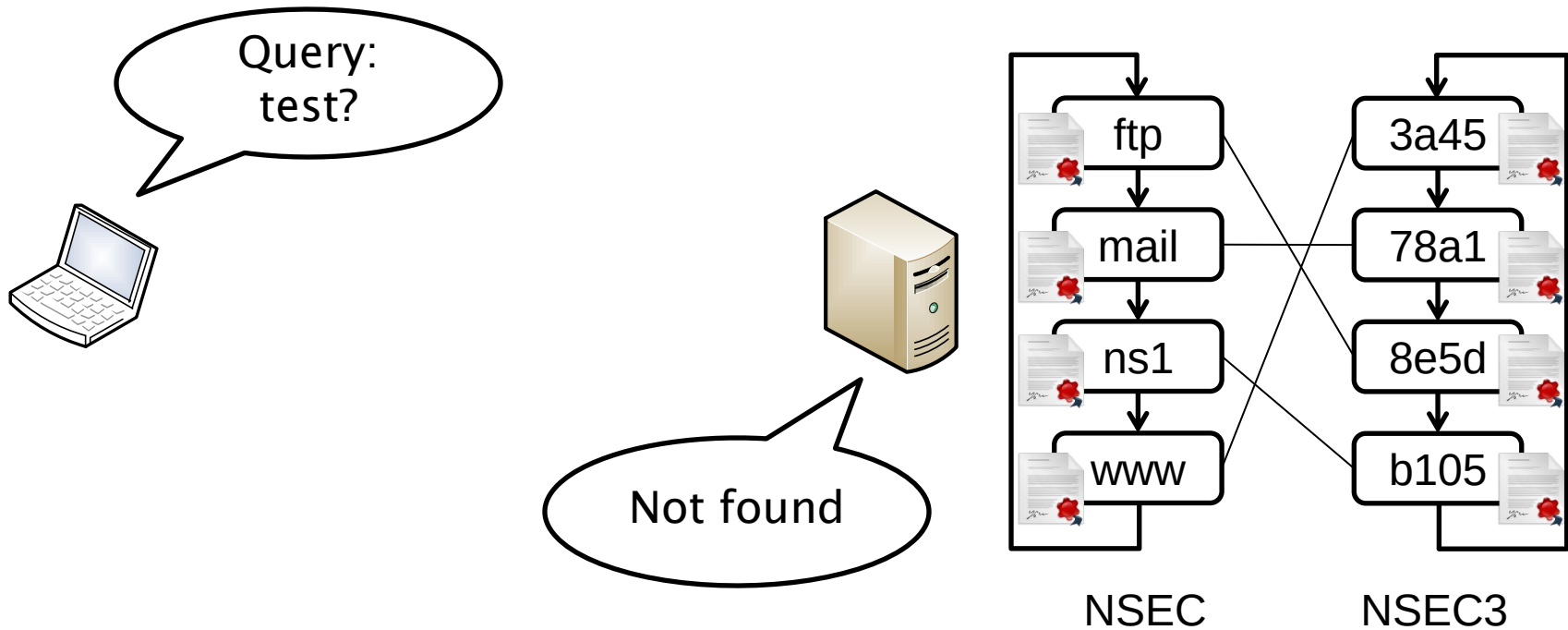
„In conflict with Germany's Federal Data Protection Act“



# Secure Denial of Existence with DNSSEC

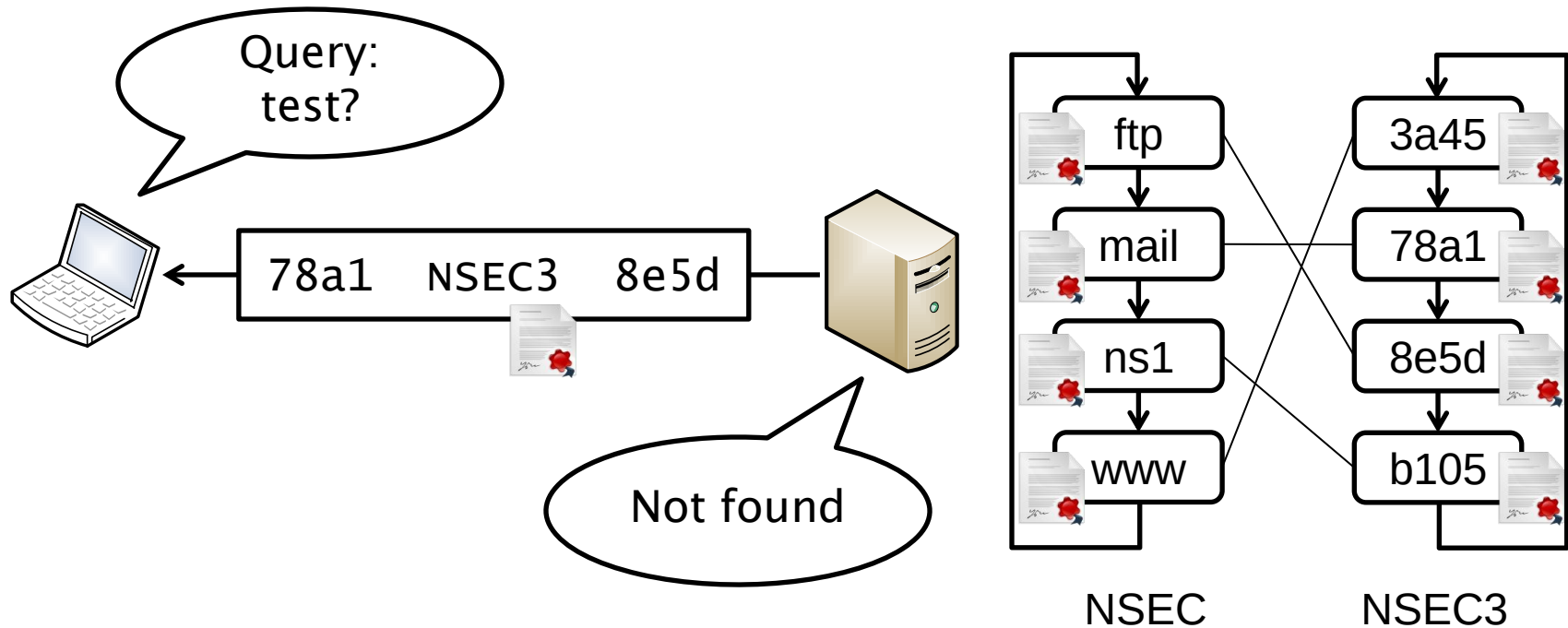


# Secure Denial of Existence with DNSSEC



- Hash query name:  $h(n)$

# Secure Denial of Existence with DNSSEC



- Hash query name:  $h(n)$
- Response:  $(h(r_1), h(r_2))$  with  $h(r_1) < h(n) < h(r_2)$

# Reasons for and against NSEC3



- Privacy
- “Opt-out” feature
- CPU load
- Message size

⇒ How well does NSEC3 prevent zone enumeration?

# NSEC3 Definition

$$\begin{aligned} h(n, s, 0) &= \text{SHA1}(\overset{\text{domain name}}{\downarrow} n \parallel \overset{\text{salt}}{\swarrow} s) \\ h(n, s, i) &= \text{SHA1}(\underset{\substack{\uparrow \\ \text{additional iterations}}}{h(n, s, i-1)} \parallel s), \text{ for } i > 0 \end{aligned}$$

- Repeated SHA1 computation
- Salt is identical for all names in a DNSSEC zone

# Top-Level Domain Survey

| TLD  | i  | salt             | TLD   | i | salt             |
|------|----|------------------|-------|---|------------------|
| at.  | 5  | b4c3ee8e123154f8 | info. | 1 | d399eaab         |
| ch.  | 2  | e7bd8151         | jp.   | 8 | f5435b71d7       |
| cn.  | 10 | aef123ab         | net.  | 0 |                  |
| com. | 0  |                  | nl.   | 5 | c9545f07a61d0d33 |
| de.  | 15 | ba5eba11         | org.  | 1 | d399eaab         |
| eu.  | 1  | 5ca1ab1e         | ru.   | 3 | 00ff             |
| fr.  | 1  | 41b69d30         | uk.   | 0 |                  |

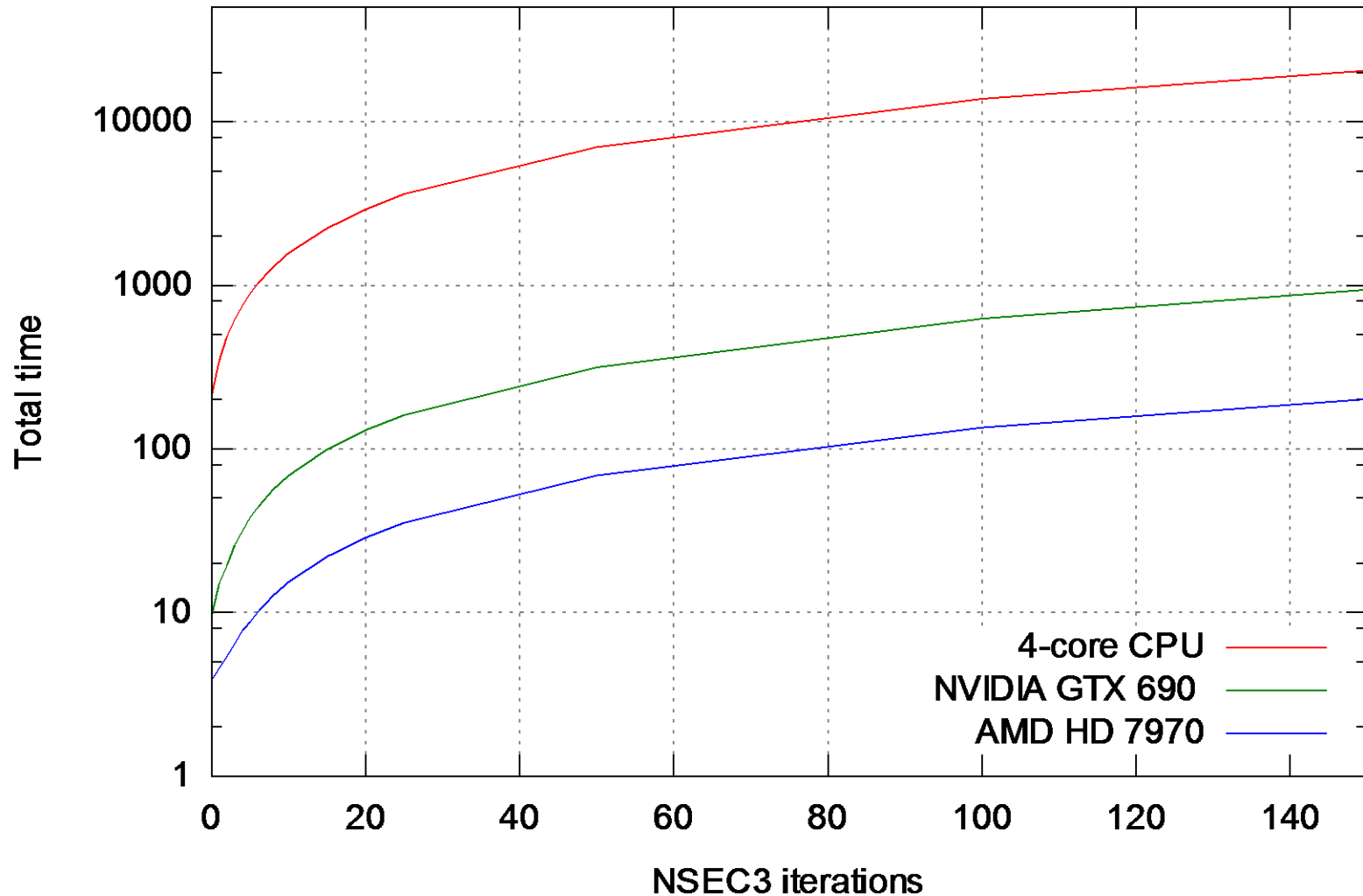
May 2014

- 335 use NSEC3, 53 use NSEC, 185 not signed
- 60% of TLDs with NSEC3 use  $i=1$
- 42% of TLDs with NSEC3 change the salt

# ATTACKING NSEC3

1. Collect hashes
2. Reverse hashes

# Why AMD Graphic Cards?



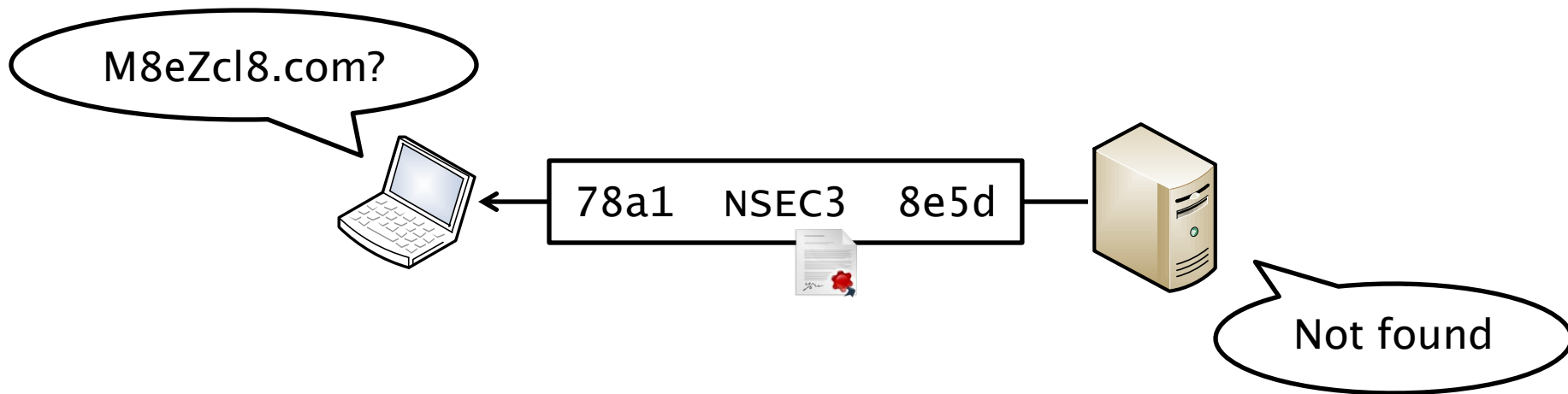


# ATTACKING NSEC3

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# Hash Crawling

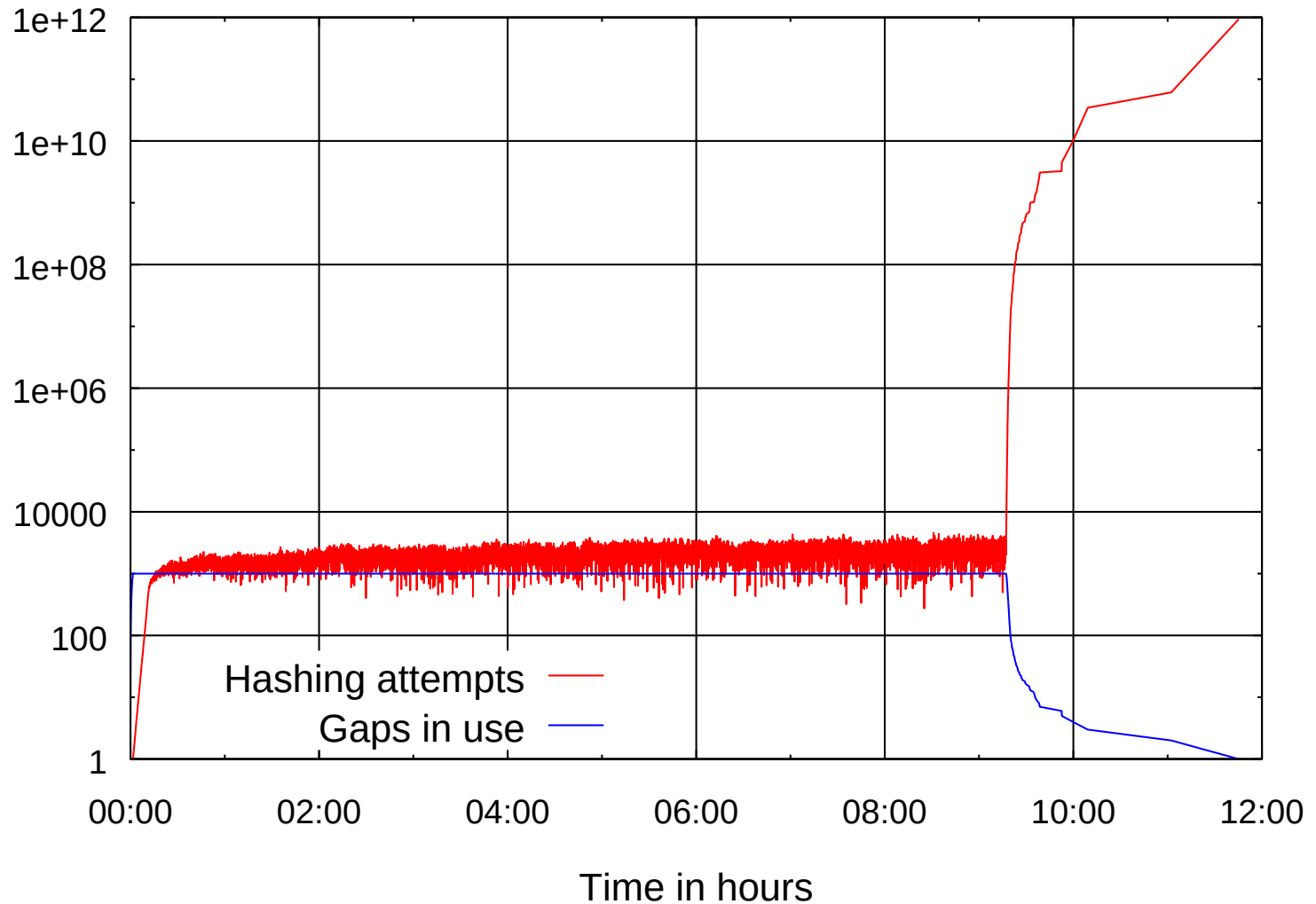
- Retrieve NSEC3 hashes from name server
  - Send queries for random non-existing names



- Check hash before sending query



# Crawling 345,000 Hashes from .com

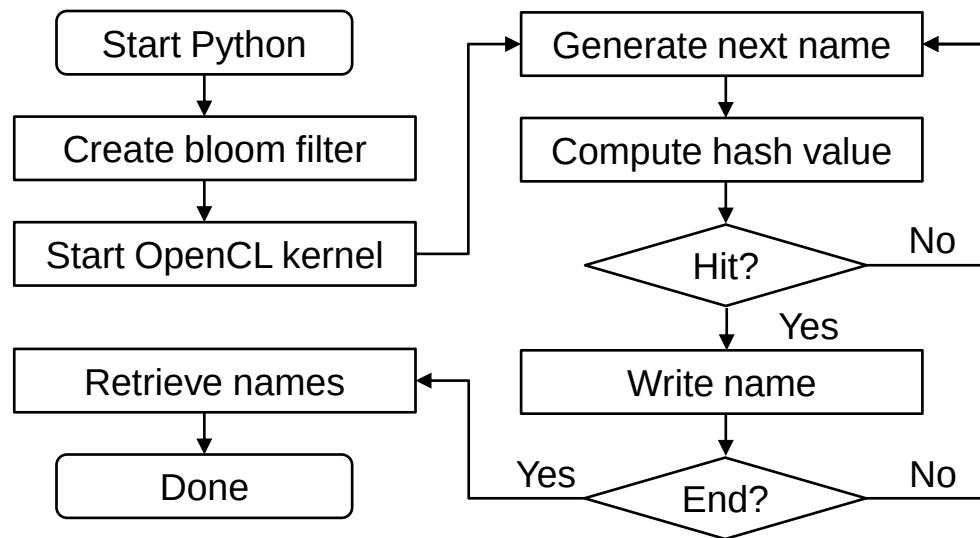


# ATTACKING NSEC3

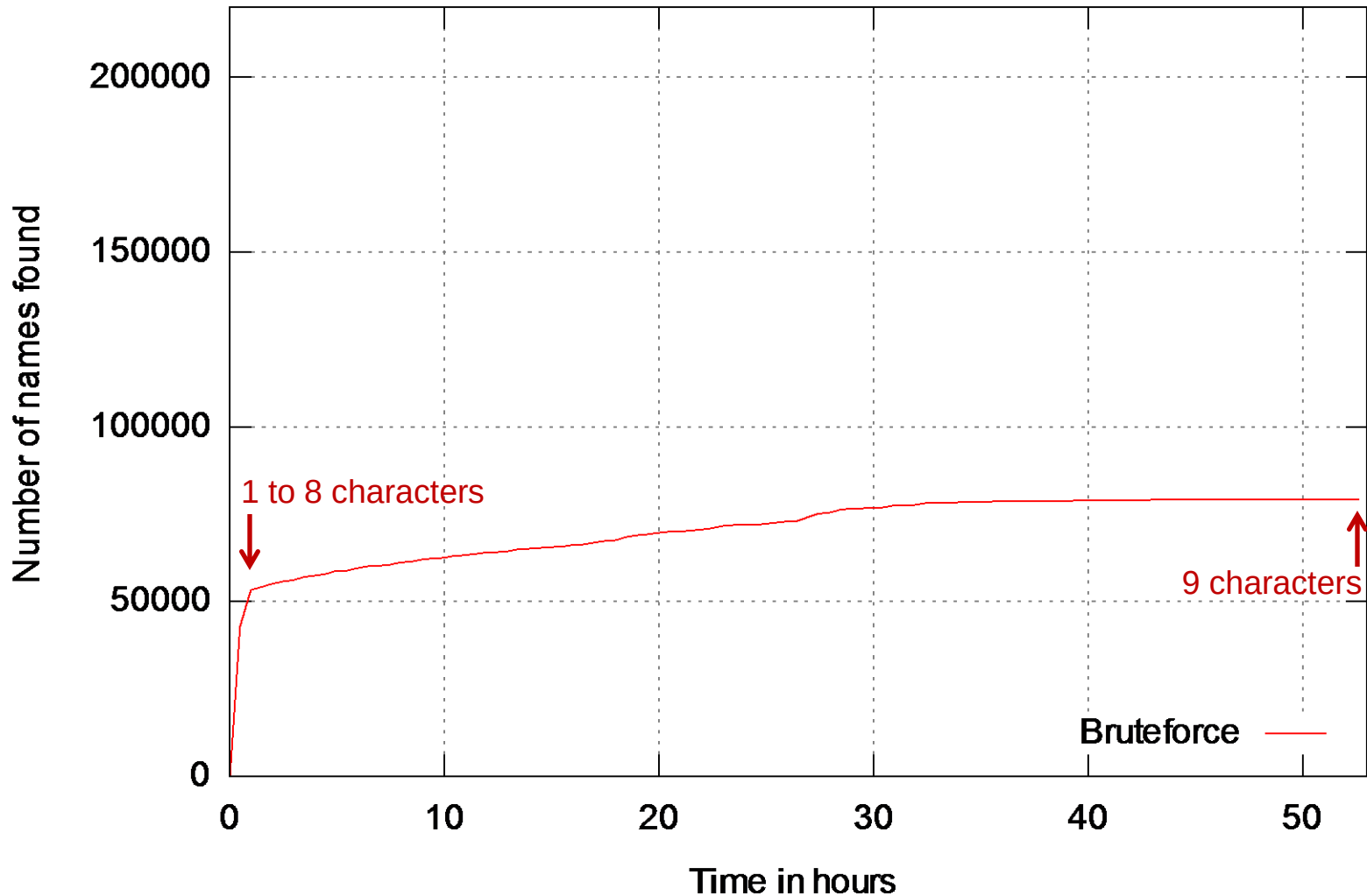
1. Collect hashes
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# GPU-based NSEC3 Hash Breaking

- Iterate through cleartext candidate names
  - Compute NSEC3 hash, compare with known hashes
- Global memory access on GPU is expensive
  - Bloom filter for 300% speedup



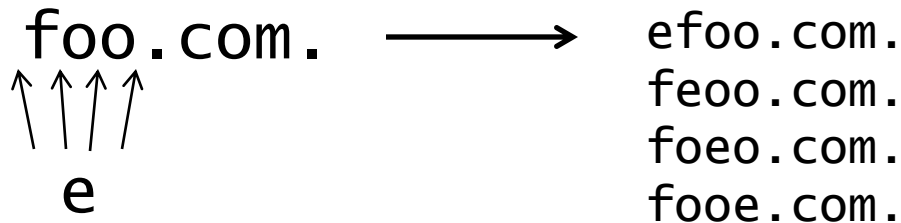
# Brute-Force Attack



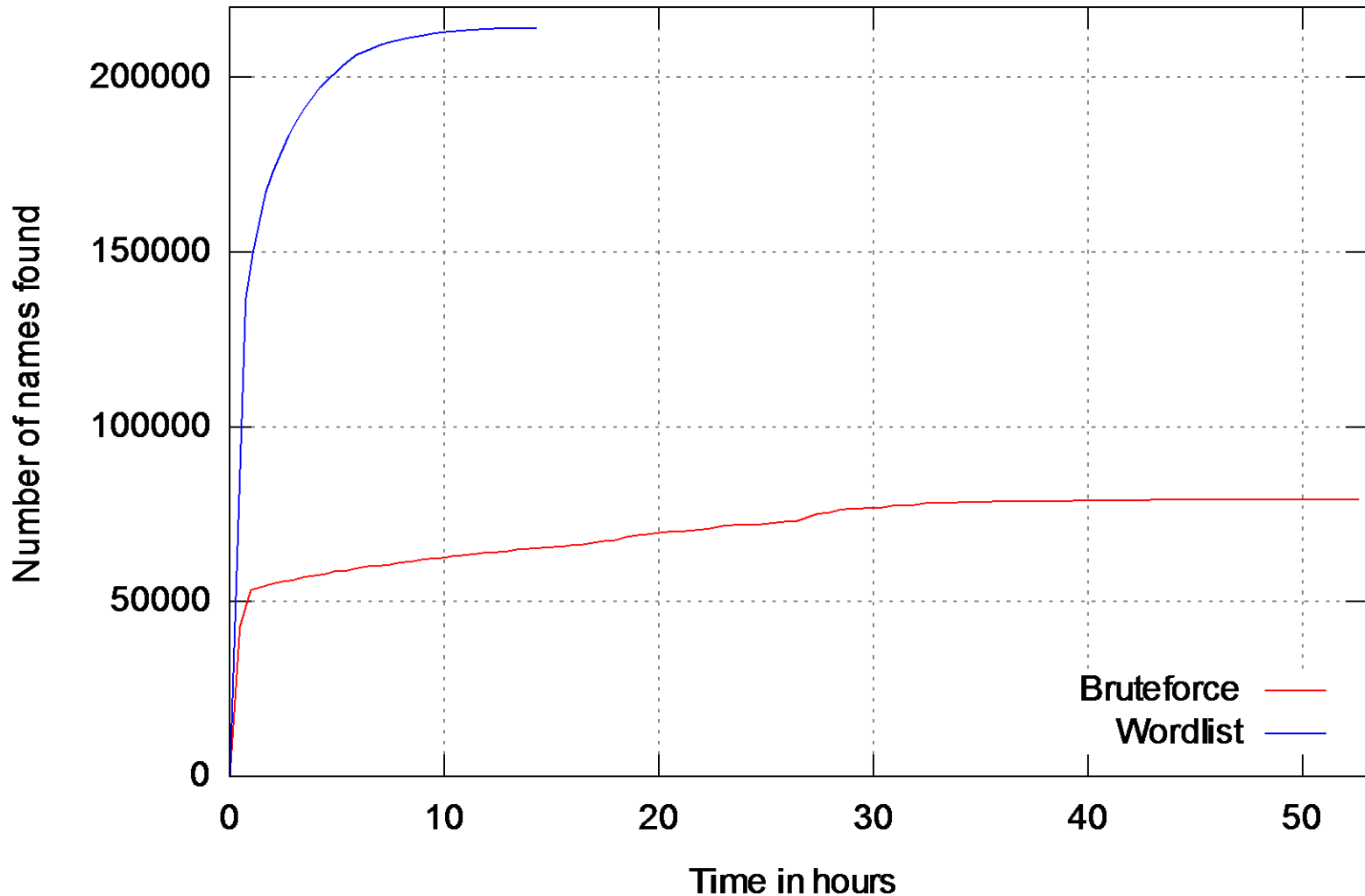
# Dictionary Attack

- Effectiveness depends on dictionary quality
  - Good public lists (Alexa, Quantcast)
  - List of all IPv4 reverse mappings (not exhaustive!)
- 7.1 mio candidate names
- Generate more candidates by inserting strings
  - Insert most common 200,000 n-grams (with  $1 \leq n \leq 15$ )

foo.com.      →      efoo.com.  
                                 feoo.com.  
                                 foeo.com.  
                                 fooe.com.

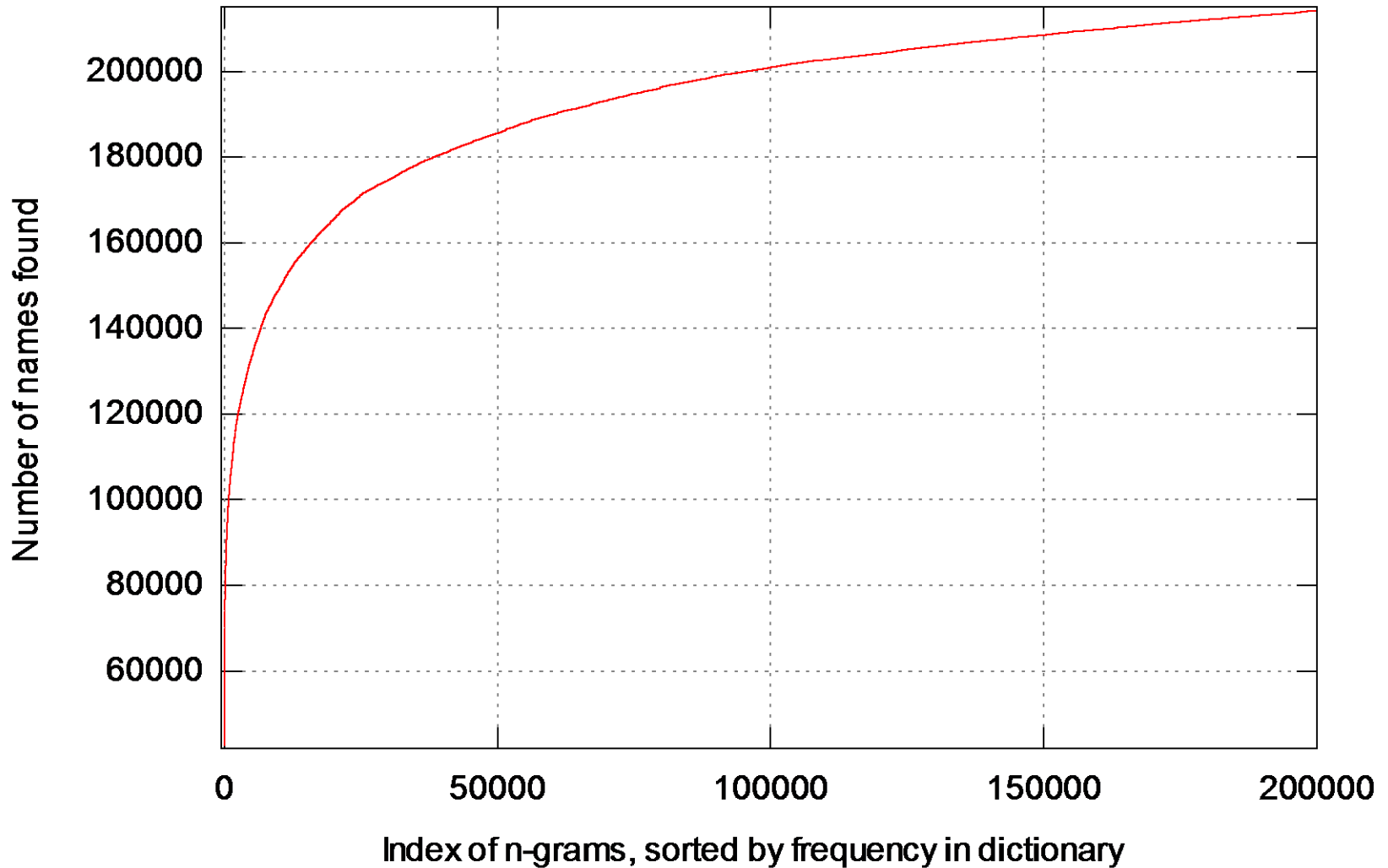


# Dictionary Attack





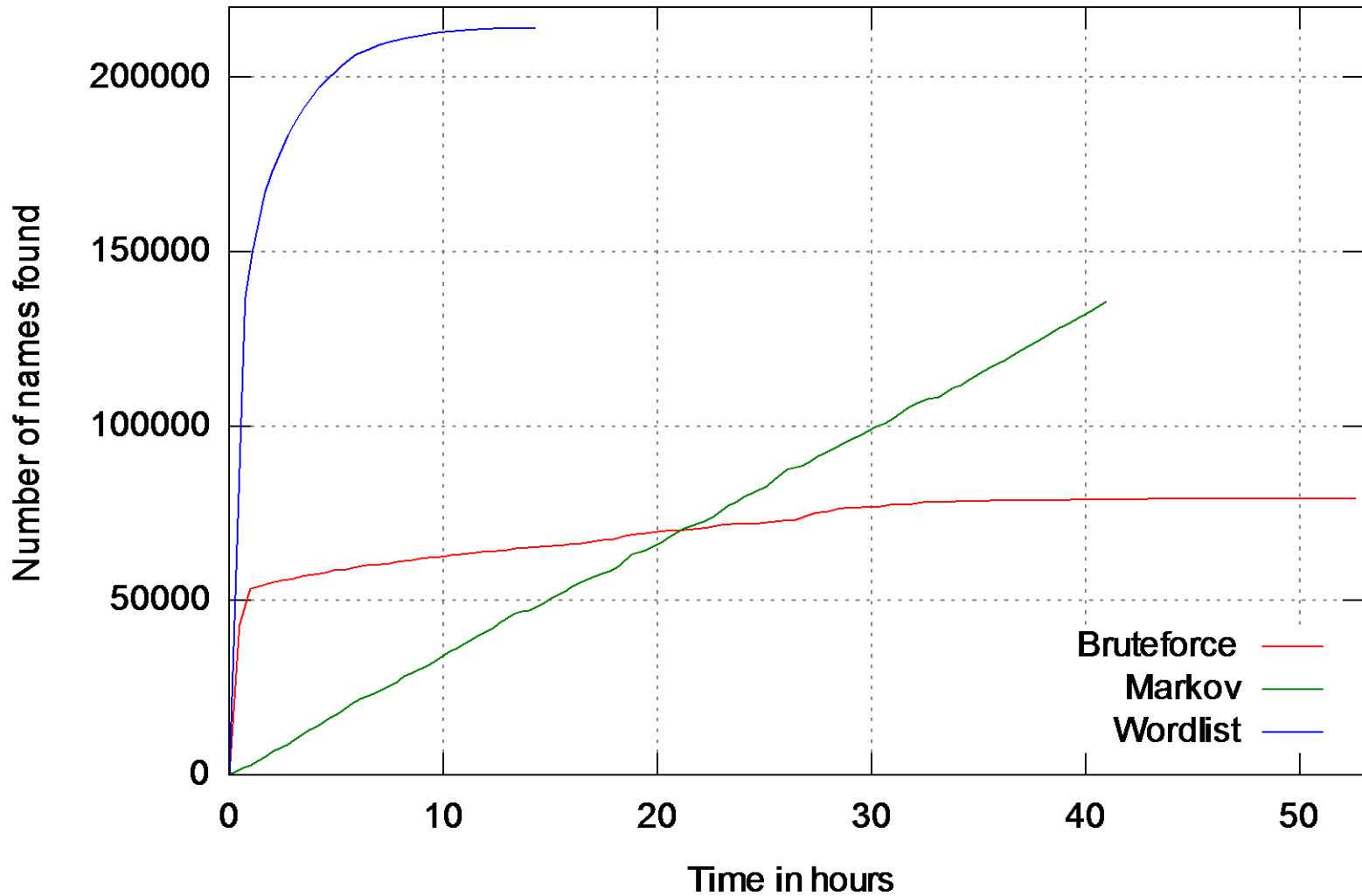
# Effectiveness of Insertion Wordlist



# Markov Attack

- Some strings are more probable than others
  - ,th‘ more common in English than ,tx‘
  - Model language with 1<sup>st</sup> order Markov chains
- Markov attack requires training
  - Use hits from brute-force and dictionary attacks
  - Markov model derived from names found
- Enumerate most probable names
  - Omit others (with given time frame)

# Markov Attack

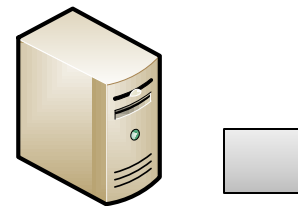


# Discussion

- Reversed 65% .com domains after 5 days
- Server operator pays for NSEC3 with CPU load
- Adjust iterations to increase CPU workload



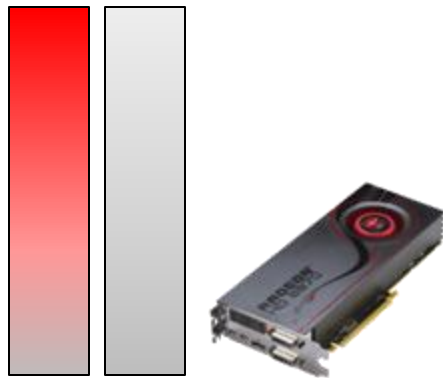
Attacker



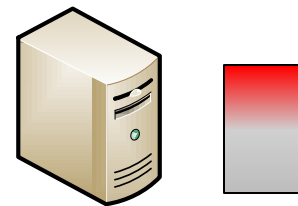
Server operator

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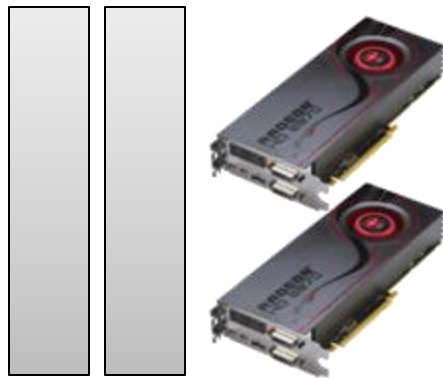
Attacker



Server operator

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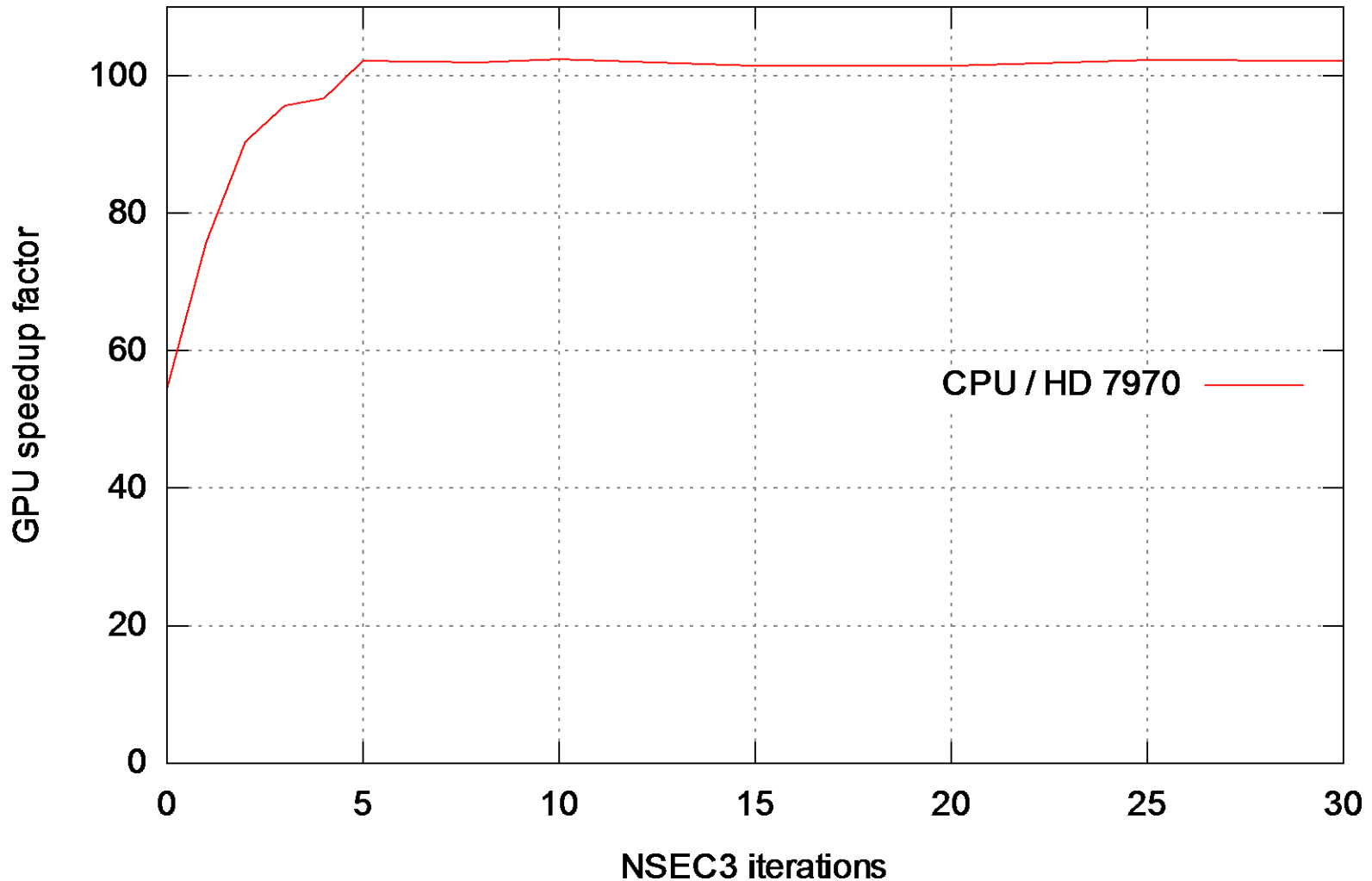


Attacker



Server operator

# Efficiency of CPU vs. GPU



# Conclusions

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- NSEC3 returns negative DNSSEC responses without disclosing existing domain names
- GPUs very efficient for breaking NSEC3 hashes
- Reversed 65% .com domains after 5 days
- Assess whether weak privacy justifies the costs
- Future work:
  - GPU-based NSEC3 accelerators for DNSSEC servers